

**Introduction to Symbolic Computation for Engineers**

**Juana Sendra, Sonia Rueda**

**(Universidad Politécnica de Madrid)**

# **3.-INTRODUCTION TO THE COMPLEXITY OF ALGORITHMS.**

- Let  $\mathcal{P}$  be the problem that we want to solve
- Let  $\mathcal{A}$  be an algorithm used to solve  $\mathcal{P}$
- $\mathcal{I}_{\mathcal{A}} = \{ \text{inputs of } \mathcal{A} \}$

We define the **COMPUTING TIME FUNCTION** as:

$$t_{\mathcal{A}} : \mathcal{I}_{\mathcal{A}} \longrightarrow \mathbb{R}$$

$$x \longrightarrow t_{\mathcal{A}}(x)$$

where  $t_{\mathcal{A}}(x)$  represents the number of **basic steps** needed for executing the algorithm  $\mathcal{A}$  to obtain the output  $\mathcal{A}(x)$ .

- In general, computing the function  $t_{\mathcal{A}}$  is not easy. For this reason we deal with simpler functions.

## Complexity Functions

We consider a partition of the set of inputs of the algorithm:

$$\mathcal{I}_A = \bigcup_{j \in J} \mathcal{I}_A(j)$$

where

- $\mathcal{I}_A(j)$  is a finite set.
- $\mathcal{I}_A(j) \cap \mathcal{I}_A(k) = \emptyset$  if  $j \neq k$ .

## Definition:

### ■ Maximum Computing Time Function

$$t_{\mathcal{A}}^+ : J \longrightarrow \mathbb{R}$$

$$j \longrightarrow t_{\mathcal{A}}^+(j) = \max\{t_{\mathcal{A}}(x) \mid x \in \mathcal{I}_{\mathcal{A}}(j)\}$$

### ■ Minimum Computing Time Function

$$t_{\mathcal{A}}^- : J \longrightarrow \mathbb{R}$$

$$j \longrightarrow t_{\mathcal{A}}^-(j) = \min\{t_{\mathcal{A}}(x) \mid x \in \mathcal{I}_{\mathcal{A}}(j)\}$$

- **Average Computing Time Function**

$$t_{\mathcal{A}}^* : J \longrightarrow \mathbb{R}$$

$$j \longrightarrow t_{\mathcal{A}}^*(j) = \sum_{x \in \mathcal{I}_{\mathcal{A}}(j)} p(x) t_{\mathcal{A}}(x)$$

where  $p(x)$  is the probability of  $x$ .

**Remarks:**

- These functions depend on the partition.
- In general we analyze  $t_{\mathcal{A}}^+$ .

## Comparing complexities

Given two algorithms  $\mathcal{A}_1$  and  $\mathcal{A}_2$  to solve the same problem  $\mathcal{P}$  we want to compare their complexity functions  $t_{\mathcal{A}_1}^+$  and  $t_{\mathcal{A}_2}^+$ .

For this purpose we introduce the following notions.

**DEFINITION:** Let  $S$  be a set and  $f, g : S \rightarrow \mathbb{R}$ .

- $f \preceq g$  ( $g$  **dominates**  $f$ ) (or  $f$  **is of order**  $g$ ) if

$$\exists c \in \mathbb{R}^+ / \forall x \in S |f(x)| \leq c|g(x)|.$$

- $f \sim g$  ( $g$  and  $f$  are **codominant**)  $f \preceq g$  and  $g \preceq f$ .

**REMARKS:**

- $f \preceq g$  can be denoted as  $f = \mathcal{O}(g)$ .
- If  $0 \leq f \leq g$  then  $f \preceq g$ .